# Results and Code for a Software-based Implementation of ECDSA for

# Control System Data Authentication and Verification

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Abstract— In previous works the need for control system data authentication and verification was discussed, and the use of the Elliptic Curve Digital Signature Algorithm (ECDSA) was proposed as a solution. Following up on that work, this paper presents a software based implementation of ECDSA utilizing ISaGRAF SoftPLCs running on both a virtual machine configuration (test string 1) and on four General Micro Systems VS 275 single board VME computers (test string 2). The implementation is primarily written in IEC 61131-3 ladder logic using specialized function blocks developed in C utilizing the OpenSSL library. The four SoftPLCs were configured in a round robin architecture in order to determine the average time it takes to generate a message, sign it, transmit it, receive it, and verify the signature. Test string 1 gave an average time of 26-28ms with PLC scan times of 22-24 ms. Test string 2 gave an average time of 55-60ms with PLC scan times of 55-60 ms. As expected, timing of the completed transmission is linearly related to PLC scan time.

Index Terms—PLC, HMI, Stuxnet, Smart Grid, NGIPS, ECDSA, Machinery Control System

### **1. INTRODUCTION**

In previous works [1] the need for increased security for control systems was discussed. The requirements for control system security were surveyed and it was determined that a method for verifying the authenticity and integrity of the commands and data transmitted between PLCs was needed. Specifically, it was proposed that the data traveling across any controls network is first signed by the sending controller at the application layer. The signature would then be transmitted across the control network along with the data and be subsequently verified by the receiving controller at the application layer. This would allow the receiving controller to verify both the authenticity and integrity of the data and would allow both controllers to utilize any available communications protocol to transmit. It was also noted that signing and verifying the data at the application layer will also allow intermediary controllers to relay data silently without additional overhead that would impact transmission time.

To accomplish this, the use of a digital signature algorithm was proposed using a public key cryptography system. A literature survey was performed of two major cryptographic technologies used in digital signature algorithms: RSA and Elliptic Curve Cryptography (ECC). Figures 1 through 3 below show timing comparisons of the two different algorithms based on the work performed by Jansma and Arrendondo. The data showed that as higher levels of security are needed ECC technology begins to outperform RSA technologies. Given that the need for security will only increase, the Elliptic Curve Digital Signature Algorithm (ECDSA) was proposed as a solution for control system data authentication and verification.

Key Length	ı		Time (s)	
ECC	R	SA	ECC	RSA
1	63	1024	0.08	0.16
2	33	2240	0.18	7.47
2	83	3072	0.27	9.80
4	09	7680	0.64	133.90
5	71 1	15360	1.44	679.06

Figure 1. ECC vs RSA Key Generation [2]

Key Leng	gth		Time (s)	
ECC		RSA	ECC	RSA
	163	1024	0.15	0.01
	233	2240	0.34	0.15
	283	3072	0.59	0.21
	409	7680	1.18	1.53
	571	15360	3.07	9.20

Figure 2. ECC vs RSA Signature Generation [2]

Key Lei	ngth		Time (s)	
ECC		RSA	ECC	RSA
	163	1024	0.23	0.01
	233	2240	0.51	0.01
	283	3072	0.86	0.01
	409	7680	1.80	0.01
	571	15360	4.53	0.03

Figure 3. ECC vs RSA Signature Verification [2]

After reviewing the mathematics of ECDSA and potential variations of the algorithm, it was concluded that ECDSA is a promising solution and that an IEC61131-3 compliant implementation would be needed in order to use existing PLC hardware and encourage widespread adoption by control engineers. Unfortunately, such an implementation would be extremely complicated and expensive to produce. Furthermore, based on the timing data presented in [2], it seemed uneconomical to commit the resources towards development of such an implementation until the feasibility of the approach could be more fully evaluated and demonstrated.

SoftPLCs provide a unique opportunity to perform this analysis by developing a prototype system that is predominately written in IEC 61131-3 code but allows the use of specialized custom function blocks written in other high level languages. This allows the development of a prototype implementation of a control system ECDSA algorithm that is able to reuse existing software libraries in order to avoid the expense of developing ladder logic cryptographic functions. The prototype implementation, being predominately written in ladder logic, can then be reused for a full IEC611-31 implementation by simply replacing the custom function blocks with IEC61131-3 Furthermore, the large number of SoftPLCs versions. already existing in both Navy and industry applications would be better suited by the prototype implementation than the full IEC 61131-3 implementation.

This paper presents the source code, timing results, and stability results for an ISaGRAF SoftPLC implementation of ECDSA for control systems. ISaGRAF SoftPLCs are currently in use on a wide range of industrial control applications around the world, and are currently employed by the Navy in two major programs: the Littoral Combat Ship (LCS) Class and the Mobile Landing Platform (MLP) Class for ship wide machinery control. The challenges and solutions uncovered when developing the implementation are discussed, and a path forward for converting the prototype into a full IEC 61131-3 compliant implementation is presented.

# 2. SYSTEM ARCHITECTURE OVERVIEW

The primary goal of the prototype was to determine the overhead the ECDSA implementation would have on control system operation. In order to have effective control, the scan times for the logic must be sufficiently low enough to approximate real time operation (less than 100 ms and ideally less than 50 ms). Additionally, for PLC to PLC communications effective control requires that you are transmitting and processing at least one set of actions every 300 ms or less (ideally every 100 ms or less). There are a number of different things that can impact both scan times and transmission times, particularly for SoftPLCs running on

a Windows OS, such as network latency and individual hardware I/O access rates.

In order to obtain averages for scan times and transmission times a simple "round robin" architecture was chosen as shown in Figure 4 below. Four SoftPLCs, labelled 1 through 4, were configured so that each processes a piece of data and then subsequently transmits it down the line, repeating the process in an indefinite loop. The data transmitted included a number of simulated signals, including BOOL, INT, REAL, DATE, and STRING values that were part of a small logic simulation routine. The details of this are discussed in Section 6 below.

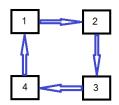


Figure 4. SoftPLC Round Robin Architecture Concept

The primary data point of interest for this system is a DINT value called COUNT. COUNT is a number that is initialized at 0 at the very start of the system. Each SoftPLC receives COUNT as an input from its predecessor, verifies the digital signature that came with the COUNT value, and if verification is successful it increments the COUNT value by 1, generates a new message, and transmits the new message to the next PLC in the loop. Figure 5 below illustrates the count increment lifecycle within the PLC's execution processes. As can be seen, each value of count represents a complete set of signature generation, data transmission, and signature verification actions. Transmission time is then synonymous with the timing of the COUNT lifecycle. In order to determine an average time, the system was run until it had reached a certain value of COUNT. The runtime was then captured and the runtime divided by the COUNT gives us the average transmission time.

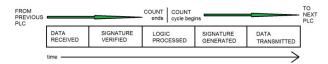


Figure 5. SoftPLC Round Robin Architecture Concept

Please note that some SoftPLC products offer an additional "WAIT" option as a step after the DATA TRANSMISSION step in their PLC. This feature normally works by predefining a PLC execution cycle time, like 100 ms, and essentially takes the 100 ms minus the time it took to complete all of the above steps and then waits for that time until it begins executing the next cycle. This is normally used to add a bit more determinism to the SoftPLC and provides a window of opportunity for the operating

system to execute tasks. Since the goal of this prototype to determine loads by monitoring scan and transmission times the WAIT feature was disabled.

The PLC scan times would be monitored by using built-in tools that come with all major SoftPLC and regular PLC packages. Scan time is not synonymous with transmission time since the PLCs are not synchronized (as is typical with industrial operations). Additionally, it is possible that a PLC might receive an invalid signature in a data transmission due to events like network transmission errors that have damaged the data. Any COUNT with an invalid signature is disregarded, making COUNT representative for only successful transmissions with correct signature generation and verification.

# **3. SOFTWARE ARCHITECTURE BACKGROUND INFORMATION**

#### **3.1. SoftPLC Package**

There are two popular SoftPLC packages currently in use by the Navy: Rockwell Canada's ISaGRAF package and Siemens WinAC. Both packages can be used to develop IEC 61131-3 ladder logic and include custom C function blocks. The primary difference between the packages is that Siemens WinAC includes additional functionality installed on the target hardware that adds real-time determinism and disables certain functions in Windows that could result in a "blue screen of death" error. ISaGRAF, on the other hand, is more comparable to Java in that it includes a simple executable known as the ISaVM that in turn executes the Like Java, ISaGRAF has multiple ISaVM logic implementations that allow the software to be used with multiple platforms including both Linux and Windows. Ultimately ISaGRAF was chosen for this implementation because it used on more US Navy Ships than Siemens WinAC and because Rockwell Canada was willing to provide free licenses to support development of this prototype.

The ISaVM executable lives with a collection of related executables in a folder that is collectively referred to as the "target". The target is installed onto the hardware via a simple copy operation which can be performed anywhere within the file structure of the SoftPLC's operating system. The target includes a main program, "ISaGRAF.exe" that is manually started by the user. This program then starts up the subprograms including both ISaVM.exe and the default Ethernet communications program "ETCP.exe". The target also includes a text file known as the "target definition" file that describes details about the target such as what custom C function blocks are available.

In order to develop custom C function blocks for inclusion in the ladder logic a tool called TDBuild (Target Definition Builder) is used to define the input and output variables of the function blocks. The tool is then used to auto-generate C code which essentially provides the API

between the ladder logic and the custom C code that was developed as part of this prototype. TDBuild is also used to update the target definition files, which are then read into the main ISaGRAF program in order to allow the software to include the specialized function blocks in the ladder logic.

For Windows-based targets the custom C code is eventually compiled as a dynamic link library (DLL) which is then copied and pasted into the target folder. This DLL, combined with the target definition files discussed previously, provide full runtime access to the custom C code for program execution. Note that you do not have to register the DLL with Windows which greatly simplifies the install.

#### 3.2. OpenSSL

As stated previously, developing a custom IEC 61131-3 implementation of ECDSA is a complicated endeavour and a prototype system that utilizes existing implementations is a necessary first step. For this prototype the OpenSSL [3] implementation has been chosen for both its efficiency and its free and open source availability to all controls engineers. There are three methods by which OpenSSL can be included in a project:

- 1. By installing OpenSSL onto the target system and using its command line interface.
- 2. By performing a fresh compile of OpenSSL as a dynamic library which can then be copied onto the target system and used after first properly registering the DLL.
- 3. By performing a fresh compile of OpenSSL as a static library which can then be included in your application without having to first install OpenSSL onto the target system or register any DLLs including the DLL used by the ISaGRAF software itself.

For this prototype OpenSSL was compiled as a static library in order to simplify the installation process onto the target machines (i.e. to keep the ISaGRAF target install down to a simple copy operation). The downside to this option is that it does complicate the structure of the actual C source code and increased development time of the prototype. It was decided that the increase in development time is justified by the increased ability to use the prototype on existing Navy ship classes that are using SoftPLCs and do not require a 100% IEC 61131-3 implementation, such as LCS and MLP.

The following steps were performed to compile OpenSSL as a static library on a Windows 7 machine:

- 1. Install Visual Studio 2010
- 2. Install Active Pearl 32-bit
- 3. Download the latest version of OpenSSL
- 4. Open the Visual Studio Command Prompt
- 5. Unzip OpenSSL to a directory such as C:\openssl-src-32
- 6. CD to the directory

- 7. Run the command: perl Configure VC-WIN32 prefix=C:\Build-OpenSSL-VC-32
- 8. Run the command: ms\do\_ms
- 9. Run the command: nmake -f ms\nt.mak
- 10. Run the command: nmake -f ms\nt.mak install

The compiled library will end up in the C:\Build-OpenSSL-VC-32 directory.

#### 3.3. Visual Studio 2010 Express

Visual Studio 2010 Express (a free product) was used to develop and compile the custom C function blocks into a DLL. Taking the C code and header files generated by TDBuild and adding it to the Visual Studio project is a simple matter, but there are additional steps that must be taken in order to configure the compiler to properly utilize the dependent libraries in order to generate the DLL. The two critical steps required are:

- 1. Under Linker General Additional Library Dependencies add the OpenSSL static library.
- 2. Under Linker Input Additional Dependencies add the libeay32.lib and the ssleay32.lib.

#### 3.4. Cryptographic Algorithms Used

As discussed in the previous work [1] performing a digital signature requires the use of two different cryptographic algorithms: a hash function and a public-key cryptography function. The hash function is performed first in order to process the arbitrarily long amount of data into a fixed length tag. The cryptographic function then encrypts the data using the private key, which can then be decrypted using the public key by an agent desiring to verify the public signature.

For this work the SHA-512 algorithm was used in combination with the P-521 ECDSA algorithm. The goal was to use the strongest measure of security possible at the time of development, and it is worth noting that the weaker P-384 ECDSA algorithm is considered valid for even TOP SECRET data. OpenSSL includes implementations of both algorithms in the static library that was compiled.

Using the SHA-512 algorithm was relatively trivial, but the P-521 algorithm proved to be a bit tricker than originally anticipated. The P-521 algorithm consists of 65 bytes of data plus 1 bit. The extra bit, when true, results in an output of 66 bytes. When the extra bit is false, the result is an output of 65 bytes. As will be illustrated in the walkthrough of the C code (and as is shown in the source code comments) special considerations were needed to monitor the output of the algorithm to ensure that the correct amount of data was read.

#### 3.5. Base64 Encoding / Decoding

Transmitting cryptographic keys and digital signatures between PLCs is not a trivial manner. PLCs only include a predefined number of data types such as BOOL, INT, DINT, and STRING which can be transmitted between devices. All of the data types except for STRING are therefore unsuitable for transmitting the keys and signatures due to the insufficient bit length.

In most PLC products, the STRING variable can consist of any ASCII character array up to a length of 255 characters according to the literature of the various vendors. This would appear to give a data value of 2040 bits using ASCII encoding, sufficiently long enough to transmit ECDSA keys and signatures. Note that this length is still too short for the use of RSA keys and digital signatures which is one of the reasons why ECC technologies were chosen.

The problem is that the 255 character set is misleading, since in ISaGRAF and in other PLC products 3 of the 255 characters are reserved to process the STRING. One of the three character slots must be used for the NULL character (which should indicate to the experienced C programmer that the STRING data type is really a C-String and not the more advanced string data type found in C++ or other high level languages). Additionally, the STRING must begin with an apostrophe and end with an apostrophe. The "apostrophe" problem becomes our greatest concern. When representing an ECDSA signature or key in ASCII it is possible that the resulting ASCII string will output an apostrophe in the middle of the key or signature. The PLC will interpret the second apostrophe (and note that you must always start with an apostrophe) as the end of the string and subsequently cut off the remaining data.

In order to work around this problem, a different character encoding must be used that maximizes the data compression of the string representation. Hexadecimal representation of the data was considered, but at only 4 bits per character it was considered insufficient. A custom variant of ASCII was considered that didn't use the apostrophe, but in the interests of conforming to widely accepted standards this idea was dropped. Base64 encoding, using 6 bits per character, became the most logical choice.

Base64 encoding allows the cryptographic keys to be transmitted as a single string variable for each key. The private key is represented by a 75 character string (plus 3 for the null and the leading/lagging apostrophes). The public key, which represents a point on an elliptic curve, is represented by a 152 character string (plus 3). The digital signature is represented by two values, SIG\_R and SIG\_S, each with a length equivalent to the private key. These values were transmitted as two separate strings to help facilitate the need for possible future expansion.

Note that the STRING variables themselves which are storing the data in the PLC ladder logic are required to be set to a fixed length. This is a fairly standard requirement by most PLC manufacturers. In order to facilitate the need for future strong cryptography it was decided to set the length for these variables in the ladder logic to the largest possible value. As a consequence, the PLC STRING variables include padded data in addition to the cryptographic keys or digital signatures.

#### 3.6. Self-Signing Keys

One of the largest areas of discussion with using a publickey cryptography system for control system security has been a concern about how to setup a public-key infrastructure (PKI). In order for two systems to be able to communicate securely, there needs to be some initial trust relationship established so that a receiving system which obtains a public key from a sending system knows that the public key really belongs to that system. Put another way, how does the receiver really know that the sender is who it claims to be?

In traditional information systems a Certificate Authority (CA) is used to issue certificates to users. This certificate contains the identity of the key pair owner, the owners public key, and a digital signature of the Certificate Authority. When users communicate securely, they receive a copy of each other's signed certificate (minus the private key) and they establish trust in each other's identity based on the successful validation of the CA's signature. Put another way, they trust each other because someone else told them it was okay. Ultimately, even trusting the digital signature of a CA is based on the idea that some human ultimately made the decision that the CA who issued the certificates is of good repute.

This kind of complicated infrastructure is necessary in IT systems, since two communicating systems may send a wide range of different kinds of data with a wide disparity in both timing and content. Furthermore, an IT system may communicate only once with another system it never heard of before and then never communicate to that system again. This constant flux requires a PKI in order to establish trust relationships with CAs and individual users. Unfortunately, running a PKI system requires a significant amount of processing power and adds significant complexity to a control system. Significant research is ongoing to develop a modified version of a PKI that can be executed on operational technology platforms, however most of these are simply slightly modified versions of the same technologies used in IT systems.

It is the opinion of the author that implementing a PKI for control systems is unnecessary. The two central problems that a PKI system solves are 1) establishing a trust relationship between two systems and 2) handling the large amount of flux in data transmitted between control systems. In Operational Technology systems, neither of these problems really exists. OT systems are designed to communicate in a very consistent and precise manner, transmitting the same basic message structure at a consistent interval to the exact same targets. The relationship between these targets is established when the system is first commissioned for operation (i.e. on the day of birth) and no new communication partnerships are ever established without significant software changes and a recommissioning of the system.

Therefore, in designing this prototype, a different system for establishing communications and updating keys was During commissioning, each system is created. preconfigured with a public / private key pair, and the public keys of each part of the system are given to each of the other parts. When an agent in the system decides to change its public / private key pair (this decision is made at regular intervals with the interval length configured at commissioning) it first generates a new key pair and then signs the new public key with the old private key. This information is then transmitted to each of the other agents in the system who then verify the key change message with the old public key they currently have on file. If the digital signature is valid, the new public key is accepted and the old public key is disregarded.

In order to ensure that an adversary will not be able to attack the system using the original preconfigured public/private key pair each agent immediately changes its key to a new random key pair at startup, before any other logic is processed. Additionally, since there is no CA and communications are highly deterministic there is little risk in changing the key pairs at a much more frequent interval. In the prototype, each of the four SoftPLC agents were configured to change their key pairs at 53 minutes, 59 minutes, 61 minutes, and 67 minutes respectively.

As a result, this system effectively generates keys 175,000 to 265,000 times more often the PKI CA systems used in IT applications and greatly weakens an adversaries capabilities to brute force crack the system. Ultimately, this allows for smaller key sizes to be used by OT systems and thus reduces the required processing power and memory required to effectively implement an ECDSA algorithm.

The major downside with this implementation is that an agent in the control system which is powered down must retain their current public/private key pair in memory and resume use of that key pair upon startup (though it can immediately change it). A agent which does not have this capability may default back to the original precommissioned key pair and will therefore be considered to be a bad actor by the other agents in the system, requiring a complete re-initialization of the entire system. Fortunately, this problem is easily solved by adding the required memory capabilities and is mitigated by the fact that these kinds of systems are highly redundant and designed to run uninterrupted for years without failures.

### **4. HARDWARE ARCHITECTURE**

Two test strings were used in the development and testing of the prototype, Test String 1 and Test String 2.

#### 4.1. Test String 1

Test String 1 used four Windows 7 virtual machines running on VMWare Player on top of a Windows 7 Pro 64bit machine. This system was primarily used for initial development and testing. The ISaGRAF software was run on the main machine in order to view the PLC code. The four virtual machines were bridged to the host NIC card which was then connected to a router. DHCP was used to assign IP addresses to both the host and virtual machines.

The main machine had the following relevant hardware specifications:

- I5-2500K processor, 4C, overclocked to 4.3 GHz, with 4x256KB of L2 Cache and 6MB of L3 Cache
- 16 GB RAM
- 1 GB Hardwired NIC
- 1.5 TB SATA 3 HDD

#### 4.2. Test String 2

Test String 2 used four General Micro System VS275 Single Board Computer VME boards in the same VME chassis. No communications across the VME backplane were used. Each board was running Windows XP Professional, 32-bit, SP2. Figure X below shows the front panel of the VS275.

Figure 6. SoftPLC Round Robin Architecture Concept

Each VS275 board had the following specifications:

- 2.16 GHz Core 2 Duo Processor, 4-MB L2 cache
- 3 GB of 667-MHz DDR-2 SDRAM
- 1 GB Hardwired NIC
- 64 GB SATA2 SSD

Hardwired Ethernet communications were used via the ENET port on the front of the board. All four boards were connected to an 8 port flat 10/100 MB hub that was in turn connected to a 10/100/1000MB router. A development station running ISaGRAF was connected to the router directly in order to download the software to the boards and to go online to the SoftPLCs. DHCP was used to assign IP addresses to each of the boards and development station.

### **5. SOURCE CODE WALKTHROUGH**

Copies of the source code developed for this prototype with comments are available as appendices to this paper, and key details and decisions made in designing the software have been presented in the previous sections. This subsection presents a high level description of each of the modules that together make up the prototype, in an effort to provide context for the source code modules.

Note that for the C code, the portions of the ISaGRAF code generated by the TDBuild tool are not available as that code is considered proprietary. Fortunately, the code not shown simply serves as the API to the PLC ladder logic and is specific to the ISaGRAF platform. Current users of

ISaGRAF will be able to use the same tools with the information provide in this document to quickly regenerate this code. Users of other SoftPLC products should be able to develop their own variations using the information provided in this paper.

#### 5.1. C CODE Walkthrough

The subsections below describe the model and C code for each of the custom function blocks that are used in the IEC61131-3 ladder logic. A fully compliant IEC61131-3 implementation will replace the C code of these blocks with a ladder logic version.

In developing the prototype, debug logic was added into each of the function blocks. Two different mechanisms for debug logic were included. The first is that each block outputs an integer status variable, with a 1 indicating successful operation and a negative number indicating a failure. The code of the blocks is organized internally in steps as shown in the source code and comments. A failed step will output a negative number equal to the step number (i.e. a fail on step 3 will output a -3 value for the status).

The second debug logic is normally turned off for full operation and was only added for development. This logic generates text files during operation that contain key status variables useful for debugging the code. This adds overhead to the program execution and creates a security risk, but is extremely valuable for debugging. In order to turn the logic on or off, the variable "ISaDEBUG" must be set true for on or false for off. This variable is located in the "debug.h" file, and a change in the status of the variable requires a recompile of the DLL. All the timing results presented in this paper are with the ISaDEBUG variable set to false.

#### 5.1.1. MsgGen

Figures 7 and 8 below define the details of the MsgGen block. In digital signature applications, the first step of an algorithm is to hash the data to be signed into one single value of fixed length, that will later be encrypted by the private key. The MsgGen block generates that hash value using the SHA-512 algorithm. The data to be signed is inputted into the B#, I#, R#, and STRING fields. The DATE and TIME fields should always be used for the current date and time stamps of the message or the entire algorithm is potentially subject to a replay attack. The block outputs the hash as a hexadecimal string.

Creating one generic message generation block that can be used in a wide range of applications is tricky, since one application could only need a few Booleans and another application could require 40 different real values. Therefore the block was structured generically in order to allow the widest range of possible inputs, and then a HASH\_I value was added that allows the user to chain multiple blocks together. Readers familiar with the SHA-512 algorithm will recognize that there is no loss in processing the data in this manner, since the algorithm generates the hash in a "chain" fashion naturally. This implementation simply extends the chain, and there is no limit to the number of MsgGen blocks that can be chained together.

The key for implementing this function is that the data to be signed must have a hash generated on both the source PLC and the destination PLC. Therefore, during system design, it is critical that the layout of the MsgGen blocks is the same on each endpoint or the signature will never verify successfully.



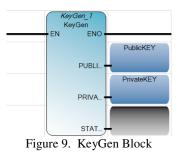
Figure 7. Three MsgGen Blocks Connected Together

Variable Name	I/O	Туре	Description
DATE	IN	DATE	Used to date the message in
			order to prevent replay
			attacks. Has resolution up
			to the second.
TIME	IN	TIME	Used to timestamp the
			message in order to prevent
			replay attacks occurring
			within the last second. Has
			resolution up to the
			millisecond.
HASH_I	IN	STRING(252)	Hash from previous
			MsgGen. Used to string
			multiple blocks together in
			order to generate messages
			containing large amounts of
			data.
B1 Thru B8	IN	BOOL	Boolean data inputs.
I1 Thru I4	IN	DINT	Integer data inputs.
R1 Thru R4	IN	REAL	Real data inputs.
STRING	IN	STRING(252)	String data inputs.
STATUS	OUT	INT	Status of the computation.
HASH_Q	OUT	STRING(252)	Hash of the message.
F	igure 8	. MsgGen Input /	Output Structure

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#### 5.1.2. KeyGen

The KeyGen block generates an ECC-521 prime public/private key pair. Ladder logic should be placed in the application to only enable the block at specific intervals when generation of a new key is desired. The new keys are outputted as base64 encoded strings. Figures 9 and 10 below define the details of the KeyGen block.



Variable Name	I/O	Туре	Description
PUBLIC_KEY	OUT	STRING(252)	ECC 521 public key
PRIVATE_KEY	OUT	STRING(252)	ECC 521 private key
STATUS	OUT	INT	Status of the computation
	10.1		~ ~ ~

Figure 10. KeyGen Input / Output Structure

#### 5.1.3. KeyVerify

The KeyVerify block is used as an added check to verify the integrity of the public/private key pair by using the private key to first sign an internal dummy hash and then using the public key to verify the signature. This block is solely used in the prototype to verify that the system is functioning correctly. As the maturity of the design improves it is believed that this block can be removed, thus reducing system overhead and improving communications results. Figures 11 and 12 below define the details of the KeyVerify block.

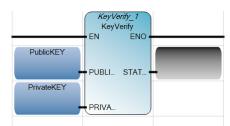


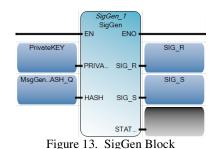
Figure 11. KeyVerify Block

Variable Name	I/O	Туре	Description
PUBLIC_KEY	IN	STRING(252)	ECC 521 public key
PRIVATE_KEY	IN	STRING(252)	ECC 521 private key
STATUS	OUT	INT	Status of the computation
Г	10 IZ		0

Figure 12. KeyVerify Input / Output Structure

#### 5.1.4. SigGen

The SigGen block is the heart of the ECDSA implementation, and takes the hash from the MsgGen block and the PrivateKey from the KeyGen block to output an ECDSA P-521 signature. The signature itself is composed of two components, SIG\_R and SIG\_S, which are each a base64 encoded string. Figures 13 and 14 below define the details of the SigGen block.



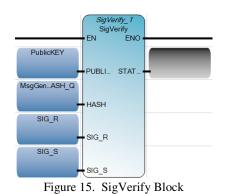
Variable Name	I/O	Туре	Description
PRIVATE_KEY	IN	STRING(252)	ECC 521 private key
HASH	IN	STRING(252)	Data hash to be signed
			from the MsgGen block.
SIG_R	OUT	STRING(252)	ECDSA P-521 signature,
			R component.
SIG_S	OUT	STRING(252)	ECDSA P-521 signature,
			S component.
STATUS	OUT	INT	Status of the computation

Figure 14. SigGen Input / Output Structure

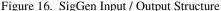
#### 5.1.5. SigVerify

The SigVerify block is used to verify that the validity of a digital signature by the receiving agent. The receiving agent first uses the MsgGen block to generate the same hash that the original sender generated, and then uses the public key of the sender along with the signature to determine if the signature and the hash match. Figures 15 and 16 below define the details of the SigGen block.

The SigVerify block is also used in the prototype system on the sender side to ensure a signature is properly generated before transmission. As the design matures it is believed that this usage of the block can be eliminated, reducing system overhead and improving communications results.



Variable Name	I/O	Туре	Description
PUBLIC_KEY	IN	STRING(252)	ECC 521 public key
HASH	IN	STRING(252)	Data hash to be verified
			from the MsgGen block.
SIG_R	IN	STRING(252)	ECDSA P-521 signature,
			R component.
SIG_S	IN	STRING(252)	ECDSA P-521 signature,
			S component.
STATUS	OUT	INT	Status of the computation
		a. a. r. i.a.	a



#### 5.2. LADDER LOGIC WALKTHROUGH

For the prototype system, the ladder logic for each of the four SoftPLC agents is structured identically, and execution of the logic commences in the following order:

- 1. ECDSA\_D#\_I process digital signatures for incoming data transmissions, including the validation of new public keys sent to the SoftPLC
- MAIN general logic such as determining current date, setting ALWAYS\_ON and ALWAYS\_OFF bits, and incrementing the counter value received during ECDSA\_D#\_I
- SIM\_BOOL and SIM\_REAL generic simulation routines that fill the place for where actual PLC logic would normally occur
- ECDSA\_KeyCntrl controls the generation and validation of the SoftPLC's own private / public key pair
- 5. ECDSA\_Q generates digital signatures for the outgoing data transmissions

#### 6. RESULTS AND FUTURE WORK

The prototype implementation was run on both test strings for periods ranging from several days to several weeks. COUNT values, reflecting the number of completed message transmissions (including both the signature generation and signature verification components) ranged from 10,000 to 4 million. In all scenarios the timing results were independent of the count length for values over 10,000. Test String 1 gave an average time of 26-28ms with PLC scan times of 22-24 ms. Test string 2 gave an average time of 55-60ms with PLC scan times of 55-60 ms. As expected, timing of the completed transmission is linearly related to PLC scan time.

Values under 10,000 counts gave slightly higher results (1-5 ms increase). During the first scan of the PLCs when the system is initializing there is an additional load on the system to establish TCP communications resulting in an increased scan time and transmission time for that cycle. For low COUNT values this initial time has a stronger effect on the average. As a result, for approximately the first 10 minutes of system runtime the average transmission time gradually decreased until the system settled at the values stated above.

#### 7. CONCLUSION

It is believed that the timing results are the worst case scenario given that no real-time modifications were added to the operating system (which would be untypical for a SoftPLC control system). Since the entire timing cycle is in the order of tens of milliseconds (as opposed to seconds as published by [2]) future development of a fully IEC 61131-3 compliant prototype is justified.

### REFERENCES

- [1] Fischer, K., "Control System Data Authentication and Verification Using Elliptic Curve Digital Signature Algorithm." Presented at ASNE Intelligent Ships Symposium X, May 22-23, 2013, Philadelphia, PA.
- [2] Jansma, N., Arrendondo, B. (2004). "Performance Comparison of Elliptic Curve and RSA Digital Signatures." (<u>http://nicj.net/files/performance\_comparison\_of\_elliptic\_curve\_and\_rsa\_digital\_signatures.pdf</u>) Accessed: 12<sup>th</sup> July 2012
- [3] OpenSSL. <u>http://www.openssl.org/</u> Accessed: 24 February 2014

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"The views expressed herein are the personal opinions of the author and are not necessarily the official views of the Department of Defense or any military department thereof."

### **APPENDIX A – C SOURCE CODE SAMPLE**

\*NOTE: Only the parts of the C Source Code written by the author are included in this Appendix. Autogenerated code from TD Build will not be included. Also note that header files not included indicate that the file is completely the result of auto-generated code. Presence of auto-generated code in the source files below will be indicated by the tag ---*autocode*---.

### A.1 - MsgGen

```
File:
                msggen.c
Author:
                kfischer
Creation date: 05/11/2013 - 19:52
POU name: MsgGen
---autocode---
//START OF CUSTOM CODE HEADERS
#include <debug.h>
#include <stdio.h>
#include <stdlib.h>
                             //for SHA512
#include <openssl/sha.h>
//END OF CUSTOM CODE HEADERS
---autocode---
  //START OF CUSTOM CODE
  //--Declarations
   /*debug variables*/
     #if ISaDEBUG
           FILE *fp;
     #endif
      /*string and hash variables*/
   #define P_PTR_HASH_I_DATA (((char*)VA_ADDRESS(pvBfData,_P_IEC_DATA->hash_i))+2)
     #define P_PTR_STRING_I_DATA (((char*)VA_ADDRESS(pvBfData,_P_IEC_DATA-
>string))+2)
      unsigned char digest[SHA512_DIGEST_LENGTH];
     char mdString[SHA512_DIGEST_LENGTH*2+1];
     char *hash_Q;
     const char *string;
     char *tstring;
     char str[STRLEN];
      char tstr[STRLEN/4];
      int i;
      /*base64 variables*/
  //--Logic
      /*STEP 0 - Initialization*/
     P_STATUS = 0;
```

```
memset(digest,0,SHA512_DIGEST_LENGTH);
memset(mdString,0,SHA512_DIGEST_LENGTH*2+1);
memset(str,0,STRLEN);
memset(tstr,0,STRLEN/4);
hash_Q = mdString;
string = str;
tstring = tstr;
/*STEP 1 - build string for hashing*/
strcpy(str, "DATE=");
sprintf(tstring,"%u,",P_DATE);
strcat(str,tstring);
strcat(str, "TIME=");
sprintf(tstring,"%u,",P_TIME);
strcat(str,tstring);
strcat(str, "HASH_I=");
sprintf(tstring,"%s,",P_PTR_HASH_I_DATA);
strcat(str,tstring);
strcat(str,"P_B1=");
sprintf(tstring,"%s,", (P_B1)?"TRUE":"FALSE");
strcat(str,tstring);
strcat(str,"P_B2=");
sprintf(tstring,"%s,", (P_B2)?"TRUE":"FALSE");
strcat(str,tstring);
strcat(str,"P_B3=");
sprintf(tstring,"%s,", (P_B3)?"TRUE":"FALSE");
strcat(str,tstring);
strcat(str,"P_B4=");
sprintf(tstring,"%s,", (P_B4)?"TRUE":"FALSE");
strcat(str,tstring);
strcat(str,"P_B5=");
sprintf(tstring, "%s,", (P_B5)?"TRUE":"FALSE");
strcat(str,tstring);
strcat(str,"P_B6=");
sprintf(tstring, "%s,", (P_B6)?"TRUE":"FALSE");
strcat(str,tstring);
strcat(str,"P_B7=");
sprintf(tstring,"%s,", (P_B7)?"TRUE":"FALSE");
strcat(str,tstring);
strcat(str,"P_B8=");
sprintf(tstring,"%s,", (P_B8)?"TRUE":"FALSE");
strcat(str,tstring);
strcat(str,"P_I1=");
sprintf(tstring,"%d,",P_I1);
strcat(str,tstring);
strcat(str,"P_I2=");
sprintf(tstring,"%d,",P_I2);
```

```
strcat(str,tstring);
strcat(str,"P_I3=");
sprintf(tstring,"%d,",P_I3);
strcat(str,tstring);
strcat(str,"P I4=");
sprintf(tstring,"%d,",P_I4);
strcat(str,tstring);
strcat(str,"P_R1=");
sprintf(tstring,"%f,",P_R1);
strcat(str,tstring);
strcat(str,"P_R2=");
sprintf(tstring,"%f,",P_R2);
strcat(str,tstring);
strcat(str,"P_R3=");
sprintf(tstring,"%f,",P_R3);
strcat(str,tstring);
strcat(str,"P_R4=");
sprintf(tstring,"%f,",P_R4);
strcat(str,tstring);
strcat(str, "P_STRING = ");
sprintf(tstring,"%s,",P_PTR_STRING_I_DATA);
strcat(str,tstring);
if (NULL == str)
{
       P_STATUS = -1; /*failed to create string*/
}
else
{
       /*STEP 2 - hash string*/
       SHA512 CTX ctx;
       SHA512 Init(&ctx);
       SHA512_Update(&ctx, string, strlen(string));
       SHA512_Final(digest, &ctx);
       for(i = 0; i < SHA512_DIGEST_LENGTH; i++)</pre>
       {
              sprintf(&mdString[i*2], "%02x", (unsigned int)digest[i]);
       }
       if (NULL == mdString)
       {
              P_STATUS = -2; /*failed to create hash*/
       }
       else
       {
              /*STEP 3 - output result*/
              strcpy(P_HASH_Q,hash_Q);
              HASH Q MAXLEN = 128;
              HASH Q CURLEN = 128;
              P STATUS = 1;
       }
```

```
/*debug logic*/
#if ISaDEBUG
fp = fopen("debug-message.txt", "w");
if (fp == NULL) {
    P_STATUS = 666;
    exit(0);
}
fprintf(fp, "string = %s\n", string);
fprintf(fp, "digest = %s\n", digest);
fprintf(fp, "digest = %s\n", ndString);

fprintf(fp, "hASH_I = %s\n", P_DATE);
fprintf(fp, "HASH_I = %s\n", (P_B1)?"TRUE":"FALSE");
fprintf(fp, "BOOL1 = %s\n", (P_B1)?"TRUE":"FALSE");
fprintf(fp, "BOOL2 = %s\n", (P_B2)?"TRUE":"FALSE");
fprintf(fp, "BOOL3 = %s\n", (P_B3)?"TRUE":"FALSE");
fprintf(fp, "BOOL4 = %s\n", (P_B3)?"TRUE":"FALSE");
fprintf(fp, "BOOL5 = %s\n", (P_B5)?"TRUE":"FALSE");
fprintf(fp, "BOOL5 = %s\n", (P_B6)?"TRUE":"FALSE");
fprintf(fp, "BOOL5 = %s\n", (P_B6)?"TRUE":"FALSE");
fprintf(fp, "BOOL5 = %s\n", (P_B8)?"TRUE":"FALSE");
fprintf(fp, "REAL1 = %i\n", P_I1);
fprintf(fp, "INT1 = %i\n", P_I3);
fprintf(fp, "REAL1 = %f\n", P_R3);
fprintf(fp, "REAL2 = %f\n", P_R4);
fclose(fp);
#endif
```

//END OF CUSTOM CODE

}

}

### A.2 – KeyGen

```
File:
                 keygen.c
Author:
                kfischer
Creation date: 15/06/2013 - 20:02
POU name:
                 KeyGen
---autocode---
//START OF CUSTOM CODE HEADERS
#include <debug.h>
#include <stdio.h>
#include <stdlib.h>
#include <openssl/ec.h> // for EC_GROUP_new_by_curve_name, EC_GROUP_free,
EC_KEY_new, EC_KEY_set_group, EC_KEY_generate_key, EC_KEY_free
#include <openssl/obj_mac.h> // for NID_secp384r1
                              // for BIGNUM conversions
#include <openssl/bn.h>
#include <openssl/bio.h>
#include <openssi/oro.../
#include <openssi/evp.h> // for base64 conversions
#include <openssl/buffer.h>
#include <math.h>
#include <base64.h>
                               // for base64 conversions
//END OF CUSTOM CODE HEADERS
---autocode---
  //START OF CUSTOM CODE
  //--Declarations
   /*debug variables*/
      #if ISaDEBUG
            FILE *fp;
            FILE *fpa;
            FILE *fpb;
            int i,j;
      #endif
  /*status variables*/
  const int set_group_success = 1;
  const int gen_success = 1;
  int set_group_status;
  int gen_status;
  /*key and base 64 variables*/
  const BIGNUM *PrivateKeyBN;
  BIGNUM *PublicKeyBN;
  const EC_POINT *PublicKeyPoint;
  EC KEY *eckey;
  EC GROUP *ecgroup;
  unsigned char *PublicKeyBinaryUnsigned, PublicKeyBinaryUnsignedArray[250];
  unsigned char *PrivateKeyBinaryUnsigned, PrivateKeyBinaryUnsignedArray[250];
```

int encodedSize;

```
char *buffer, BufferArray[250];
   unsigned char *PrivateKeyBase64, PrivateKeyBase64Array[250];
   int encodedSize b;
   char *buffer_b, BufferArray_b[250];
   unsigned char *PublicKeyBase64, PublicKeyBase64Array[250];
  //--Logic
   /*STEP 0 - Initialization*/
      P_STATUS = 0;
      memset(PrivateKeyBinaryUnsignedArray,0,250);
      PrivateKeyBinaryUnsigned = PrivateKeyBinaryUnsignedArray;
      memset(PublicKeyBinaryUnsignedArray,0,250);
      PublicKeyBinaryUnsigned = PublicKeyBinaryUnsignedArray;
      memset(PrivateKeyBase64Array,0,250);
      PrivateKeyBase64 = PrivateKeyBase64Array;
      memset(PublicKeyBase64Array,0,250);
      PublicKeyBase64 = PublicKeyBase64Array;
      memset(BufferArray,0,250);
      buffer = BufferArray;
      memset(BufferArray b,0,250);
      buffer_b = BufferArray_b;
       /*STEP 1 - create key object*/
    eckey=EC_KEY_new();
    if (NULL == eckey)
    {
        P_STATUS = -1; /*failed to create key object*/
    }
   else
    {
              /*STEP 2 - create EC GROUP object*/
        //ecgroup = EC_GROUP_new_by_curve_name(NID_secp384r1);
                                                                --worked for 384
ECDSA, will want to come back and add an option to turn this on and off
              ecgroup = EC_GROUP_new_by_curve_name(NID_secp521r1);
        if (NULL == ecgroup)
        {
            P_STATUS = -2; /*failed to create new EC Group*/
        }
        else
        {
                    /*STEP 3 - associate key with group*/
            set_group_status = EC_KEY_set_group(eckey,ecgroup);
            if (set_group_success != set_group_status)
            {
                           P_STATUS = -3; /*failed to associate key with group*/
            }
            else
                     {
                           /*STEP 4 - create PRIVATE and PUBLIC keys*/
                gen_status = EC_KEY_generate_key(eckey);
                if (gen_success != gen_status)
```

{ P\_STATUS = -4; /\*failed to generate EC Key\*/ } else { /\*STEP 5 - extract PRIVATE keys\*/ PrivateKeyBN = EC\_KEY\_get0\_private\_key(eckey); if (PrivateKeyBN == NULL) { P STATUS = -5; /\*failed to extract PRIVATE key\*/ } else { /\*STEP 6 - extract PUBLIC KEY\*/ PublicKeyPoint = EC\_KEY\_get0\_public\_key(eckey); if (PublicKeyPoint == NULL) { P\_STATUS = -6; /\*failed to extract PUBLIC key\*/ } else { /\*STEP 7a - Private Key convert from BN to binary and then encode as base64\*/ BN bn2bin(PrivateKeyBN,PrivateKeyBinaryUnsigned); encodedSize = EVP EncodeBlock(PrivateKeyBase64,PrivateKeyBinaryUnsigned,89); /\*STEP 7b - Public Key convert from Point to BN to binary and then encode as base64\*/ PublicKeyBN = BN new(); EC POINT point2bn(ecgroup,PublicKeyPoint,POINT CONVERSION UNCOMPRESSED,PublicKe yBN,NULL); BN\_bn2bin(PublicKeyBN,PublicKeyBinaryUnsigned); encodedSize\_b = EVP\_EncodeBlock(PublicKeyBase64,PublicKeyBinaryUnsigned,179); /\*STEP 7c - set outputs\*/ buffer = (char\*) PrivateKeyBase64; buffer\_b = (char\*) PublicKeyBase64; strcpy(P\_PRIVATE\_KEY\_BUFADD, buffer); PRIVATE\_KEY\_MAXLEN = 250; PRIVATE\_KEY\_CURLEN = encodedSize; strcpy(P PUBLIC KEY BUFADD, buffer b); PUBLIC KEY MAXLEN = 250; PUBLIC KEY CURLEN = encodedSize b;

```
P_STATUS = 1; /*success*/
```

```
/*debug logic*/
#if ISaDEBUG
       fpb = fopen("debug-keygen-BIGNUM.txt", "w");
       fprintf(fpb, "PublicKeyBigNumber = ");
       BN_print_fp(fpb, PublicKeyBN);
       fprintf(fpb, " \n");
fprintf(fpb, " \n");
fprintf(fpb, " \n");
fprintf(fpb, "PrivateKeyBigNumber = ");
       BN_print_fp(fpb, PrivateKeyBN);
       fclose(fpb);
       fp = fopen("debug-keygen-Private.txt", "w");
       fprintf(fp, "hello world\n");
fprintf(fp,"encodedSize = %i\n", encodedSize);
       fprintf(fp,"encodedSize_b = %i\n", encodedSize_b);
       fprintf(fp,"array dump = ");
       for (i = 0; i<250; i++)</pre>
       {
               fprintf(fp, "%c", PrivateKeyBinaryUnsignedArray[i]);
       }
       fprintf(fp, " \n");
fprintf(fp, "PrivateKeyBinaryUnsigned = %s\n", PrivateKeyBinaryUnsigned);
       fprintf(fp,"PrivateKeyBase64 = %s\n", PrivateKeyBase64);
fprintf(fp,"buffer = %s\n", buffer);
       fclose(fp);
#endif
                                              }
                                      }
                              }
                       }
               }
       }
   //END OF CUSTOM CODE
}
```

### A.3 – KeyVerify

```
File:keyverify.cAuthor:kfischer
Author:kfischerCreation date:15/06/2013 - 13:53POU name:KeyVerify
---autocode---
//START OF CUSTOM CODE HEADERS
#include <debug.h>
#include <stdio.h>
#include <stdlib.h>
#include <openssl/ec.h> // for EC_GROUP_new_by_curve_name, EC_GROUP_free,
EC_KEY_new, EC_KEY_set_group, EC_KEY_generate_key, EC_KEY_free
#include <openssl/ecdsa.h> // for ECDSA_do_sign, ECDSA_do_verify
#include <openssl/obj_mac.h> // for NID_secp384r1
#include <openssl/bn.h>
                              // for BIGNUM conversions
#include <openssl/bio.h>
#include <openssl/evp.h>
#include <openssl/buffer.h>
#include <math.h>
#include <base64.h>
                                // for base64 conversions
//END OF CUSTOM CODE HEADERS
---autocode---
  //START OF CUSTOM CODE
  //--Declarations
   /*debug variables*/
      #if ISaDEBUG
      FILE *fp;
      FILE *fpa;
      FILE *fpb;
      int i, j;
      #endif
   /*status variables*/
      #define P PTR PUBLIC KEY DATA (((char*)VA ADDRESS(pvBfData, P IEC DATA-
>public key))+2)
      #define P PTR PRIVATE KEY DATA (((char*)VA ADDRESS(pvBfData, P IEC DATA-
>private_key))+2)
      const int set_group_success = 1;
      const int gen_success = 1;
      const int set_public_success = 1;
      const int set_private_success = 1;
      const int verify_success = 1;
      int set_group_status;
      int set_public_status;
      int set_private_status;
```

```
int verify status;
   unsigned char hash[] = "c7fbca202a95a570285e3d700eb04ca2";
/*base64 variables*/
   EC KEY *eckey;
   EC GROUP *ecgroup;
   EC POINT *PublicKeyPoint;
   BIGNUM *PrivateKeyBN;
   BIGNUM *PublicKeyBN;
   ECDSA_SIG *signature;
   unsigned char *buffer;
   unsigned char *buffer_b;
   unsigned char *PublicKeyBinaryUnsigned, PublicKeyBinaryUnsignedArray[250];
   unsigned char *PrivateKeyBinaryUnsigned, PrivateKeyBinaryUnsignedArray[250];
   const unsigned char *PublicKeyBinaryConst;
   const unsigned char *PrivateKeyBinaryConst;
//--Logic
   P_STATUS = 0;
   memset(PrivateKeyBinaryUnsignedArray,0,250);
   PrivateKeyBinaryUnsigned = PrivateKeyBinaryUnsignedArray;
   memset(PublicKeyBinaryUnsignedArray,0,250);
   PublicKeyBinaryUnsigned = PublicKeyBinaryUnsignedArray;
   /*STEP 1 - create key object*/
   eckey=EC_KEY_new();
if (NULL == eckey)
{
    P STATUS = -1; /*failed to create key object*/
}
else
 {
           /*STEP 2 - create EC_GROUP object*/
          ecgroup = EC_GROUP_new_by_curve_name(NID_secp521r1);
    if (NULL == ecgroup)
     {
         P_STATUS = -2; /*failed to create new EC Group*/
    }
    else
     {
                 /*STEP 3 - associate key with group*/
         set_group_status = EC_KEY_set_group(eckey,ecgroup);
         if (set_group_success != set_group_status)
         {
                        P_STATUS = -3; /*failed to associate key with group*/
         }
         else
                  {
                        /*STEP 4 - set public key*/
                        PublicKeyPoint = EC_POINT_new(ecgroup);
```

```
PublicKeyBN = BN_new();
buffer_b = (unsigned char*) P_PTR_PUBLIC_KEY_DATA;
EVP_DecodeBlock(PublicKeyBinaryUnsigned,buffer_b,240);
PublicKeyBinaryUnsigned;
PublicKeyBinaryConst = (const unsigned char*)
BN bin2bn(PublicKeyBinaryConst,133,PublicKeyBN);
```

EC\_POINT\_bn2point(ecgroup,PublicKeyBN,PublicKeyPoint,NULL);

```
set_public_status =
EC_KEY_set_public_key(eckey,PublicKeyPoint);
    if (set_public_success != set_public_status)
    {
        P_STATUS = -4; /*failed to set public key*/
    }
    else
    {
        /*STEP 5 - set private key*/
        PrivateKeyBN = BN_new();
        buffer = (unsigned char*) P_PTR_PRIVATE_KEY_DATA;
    }
}
```

PrivateKeyBinaryUnsigned;

/\*521 ECDSA is unique due to the extra bit resulting in a length of 65.125 bytes, which means that when converting from bin to bn we will alternatively be converting 66 or 65 bytes depending on if the extra bit is a 1 or not. The presence of the extra bit is encoded in the in the first base64 characters. The byte 0000 0001 will be broken in base64 to a pair of 0000000 and 1XXXXXX, which will result in the key having a leading A character representing the 0000000. However, the byte 0000 0000 will not be encoded at all, which will result in the key not having a leading A. Detection of the leading A is therefore critical to properly convert the bin to BN.\*/

```
/*STEP 6 - create signature on dummy hash*/
                                              signature = ECDSA_do_sign(hash,32,eckey);
                                              if (NULL == signature)
                                              {
                                                      P_STATUS = -6; /*failed to generate
signature*/
                                              }
                                              else
                                              {
                                                      /*STEP 7 - verify signature*/
                                                      verify_status =
ECDSA_do_verify(hash,32,signature,eckey);
                                                      if (verify_success != verify_status)
                                                      {
                                                             P_STATUS = -7; /*verification
failed*/
                                                      }
                                                      else
                                                      {
                                                             P_STATUS = 1; /*keys verified*/
                                                      }
                                              }
                                      }
                              }
                       }
               }
       }
       /*CLEANUP - Release memory structures to prevent memory leaks*/
       EC KEY free(eckey);
                                                      /*frees EC KEY memory allocation*/
       EC_GROUP_clear_free(ecgroup);
                                           /*destroys any sensitive data in the group and
then frees the memory*/
       EC_POINT_free(PublicKeyPoint); /*frees EC_POINT memory allocation*/
       BN_clear_free(PrivateKeyBN);
                                           /*overwrites the BN before returning memory to
the system*/
       BN clear free(PublicKeyBN);
       ECDSA SIG free(signature);
                                              /*frees ECDSA SIG memory allocation*/
/*debug logic*/
#if ISaDEBUG
       fpa = fopen("debug-keyverify-public.txt", "w");
       fprintf(fpa, "hello world \n");
fprintf(fpa, "P_PTR_PUBLIC_KEY_DATA = %s\n", P_PTR_PUBLIC_KEY_DATA);
fprintf(fpa, " \n");
fprintf(fpa, " \n");
fprintf(fpa, "Public Key Array dump = ");
       for (j = 0; j<250; j++)</pre>
       {
               fprintf(fpa, "%c", PublicKeyBinaryUnsignedArray[j]);
       }
       fprintf(fpa, " \n");
fprintf(fpa, " \n");
       fclose(fpa);
       fp = fopen("debug-keyverify-private.txt", "w");
       fprintf(fp, "hello world\n");
       fprintf(fp, "P_PTR_PRIVATE_KEY_DATA = %s\n", P_PTR_PRIVATE KEY_DATA);
       fprintf(fp, "P_PTR_PRIVATE_KEY_DATA[0] = %c\n", P_PTR_PRIVATE_KEY_DATA[0]);
```

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```
fprintf(fp, "Private Key Array dump = ");
for (i = 0; i<250; i++)
{
     fprintf(fp, "%c", PrivateKeyBinaryUnsignedArray[i]);
}
fprintf(fp, " \n");
fclose(fp);

fpb = fopen("debug-keyverify-BIGNUM.txt", "w");
fprintf(fpb, "PublicKeyBigNumber = ");
BN_print_fp(fpb, PublicKeyBN);
fprintf(fpb, " \n");
fprintf(fpb, " \n");
fprintf(fpb, "PrivateKeyBigNumber = ");
BN_print_fp(fpb, PrivateKeyBN);
fprintf(fpb, " \n");
fprintf(fp
```

//END OF CUSTOM CODE

}

### A.4 – SigGen

```
File:
                 siggen.c
Author:kfischerCreation date:10/07/2013 - 22:27POU name:SigGen
---autocode---
//START OF CUSTOM CODE HEADERS
#include <debug.h>
#include <stdio.h>
#include <stdlib.h>
#include <openssl/ec.h> // for EC_GROUP_new_by_curve_name, EC_GROUP_free,
EC_KEY_new, EC_KEY_set_group, EC_KEY_generate_key, EC_KEY_free
#include <openssl/ecdsa.h> // for ECDSA_do_sign, ECDSA_do_verify
#include <openssl/obj_mac.h> // for NID_secp384r1
#include <openssl/bn.h>
                              // for BIGNUM conversions
#include <openssl/bio.h>
#include <openssl/evp.h>
#include <openssl/buffer.h>
#include <math.h>
#include <base64.h>
                                // for base64 conversions
//END OF CUSTOM CODE HEADERS
---autocode---
  //START OF CUSTOM CODE
  //--Declarations
   /*debug variables*/
      #if ISaDEBUG
      FILE *fp;
      int i;
      #endif
   /*status variables*/
      #define P_PTR_PRIVATE_KEY_DATA (((char*)VA_ADDRESS(pvBfData,_P_IEC_DATA-
>private key))+2)
      #define P PTR HASH DATA (((char*)VA ADDRESS(pvBfData, P IEC DATA->hash))+2)
      const int set group success = 1;
      const int gen_success = 1;
      const int set public success = 1;
      const int set_private_success = 1;
      const int verify_success = 1;
      int set_group_status;
      int set_private_status;
      unsigned char hash[128];
```

```
/*EC Variables*/
   EC KEY *eckey;
   EC_GROUP *ecgroup;
   ECDSA SIG *signature;
/*base64 variables for decomposing private key*/
   BIGNUM *PrivateKeyBN;
   unsigned char *buffer;
 unsigned char *PrivateKeyBinaryUnsigned, PrivateKeyBinaryUnsignedArray[250];
   const unsigned char *PrivateKeyBinaryConst;
/*base64 variables for encoding SIG_R and SIG_S*/
   int SIG R EncodedSize;
   unsigned char *SIG_R_BinaryUnsigned, SIG_R_BinaryUnsignedArray[250];
 unsigned char *SIG R Base64, SIG R Base64Array[250];
   int SIG S EncodedSize;
   unsigned char *SIG_S_BinaryUnsigned, SIG_S_BinaryUnsignedArray[250];
unsigned char *SIG S Base64, SIG S Base64Array[250];
   char *SIG_R_Buffer, SIG_R_BufferArray[250];
   char *SIG_S_Buffer, SIG_S_BufferArray[250];
//--Logic
    /*STEP 0 - Initialization*/
   P STATUS = 0;
   memcpy(hash,P_PTR_HASH_DATA,128);
   memset(PrivateKeyBinaryUnsignedArray,0,250);
   PrivateKeyBinaryUnsigned = PrivateKeyBinaryUnsignedArray;
   memset(SIG R BinaryUnsignedArray,0,250);
   SIG R BinaryUnsigned = SIG R BinaryUnsignedArray;
   memset(SIG R Base64Array,0,250);
   SIG R Base64 = SIG R Base64Array;
   memset(SIG R BufferArray,0,250);
   SIG R Buffer = SIG R BufferArray;
   memset(SIG_S_BinaryUnsignedArray,0,250);
   SIG_S_BinaryUnsigned = SIG_S_BinaryUnsignedArray;
   memset(SIG_S_Base64Array,0,250);
   SIG_S_Base64 = SIG_S_Base64Array;
   memset(SIG_S_BufferArray,0,250);
   SIG_S_Buffer = SIG_S_BufferArray;
   /*STEP 1 - create key object*/
   eckey=EC_KEY_new();
if (NULL == eckey)
 {
     P_STATUS = -1; /*failed to create key object*/
}
else
 {
           /*STEP 2 - create EC GROUP object*/
          ecgroup = EC_GROUP_new_by_curve_name(NID_secp521r1);
     if (NULL == ecgroup)
     {
```

```
P STATUS = -2; /*failed to create new EC Group*/
}
else
{
            /*STEP 3 - associate key with group*/
    set_group_status = EC_KEY_set_group(eckey,ecgroup);
    if (set_group_success != set_group_status)
    {
                    P STATUS = -3; /*failed to associate key with group*/
    }
    else
             {
                    /*STEP 4 - set private key*/
                    PrivateKeyBN = BN_new();
                    buffer = (unsigned char*) P_PTR_PRIVATE_KEY_DATA;
                    EVP_DecodeBlock(PrivateKeyBinaryUnsigned, buffer, 120);
                    PrivateKeyBinaryConst = (const unsigned char*)
```

PrivateKeyBinaryUnsigned;

/\*521 ECDSA is unique due to the extra bit resulting in a length of 65.125 bytes, which means that when converting from bin to bn we will alternatively be converting 66 or 65 bytes depending on if the extra bit is a 1 or not. The presence of the extra bit is encoded in the in the first base64 characters. The byte 0000 0001 will be broken in base64 to a pair of 0000000 and 1XXXXXX, which will result in the key having a leading A character representing the 0000000. However, the byte 0000 0000 will not be encoded at all, which will result in the key not having a leading A. Detection of the leading A is therefore critical to properly convert the bin to BN.\*/

```
if (P PTR PRIVATE KEY DATA[0] == 'A')
                            {
                                   BN_bin2bn(PrivateKeyBinaryConst,66,PrivateKeyBN);
                            }
                            else
                            {
                                   BN bin2bn(PrivateKeyBinaryConst,65,PrivateKeyBN);
                            }
                            set_private_status =
EC_KEY_set_private_key(eckey,PrivateKeyBN);
                            if (set_private_success != set_private_status)
                            {
                                   P_STATUS = -4; /*failed to set private key*/
                            }
                            else
                            {
                                   /*STEP 5 - create signature on dummy hash*/
                                   signature = ECDSA_do_sign(hash,128,eckey);
                                   if (NULL == signature)
                                   {
                                          P_STATUS = -5; /*failed to generate
signature*/
                                   }
                                   else
                                   {
                                          /*STEP 6 - base64 encode SIG R and SIG S and
output result*/
```

BN\_bn2bin(signature->r,SIG\_R\_BinaryUnsigned); SIG\_R\_EncodedSize = EVP\_EncodeBlock(SIG\_R\_Base64,SIG\_R\_BinaryUnsigned,89);

```
BN_bn2bin(signature->s,SIG_S_BinaryUnsigned);
                                          SIG S EncodedSize =
EVP_EncodeBlock(SIG_S_Base64,SIG_S_BinaryUnsigned,89);
                                          SIG R Buffer = (char *) SIG R Base64;
                                          SIG S Buffer = (char *) SIG S Base64;
                                          strcpy(P_SIG_R_BUFADD,SIG_R_Buffer);
                                          SIG_R_MAXLEN = 250;
                                          SIG_R_CURLEN = SIG_R_EncodedSize;
                                          strcpy(P_SIG_S_BUFADD,SIG_S_Buffer);
                                          SIG_S_MAXLEN = 250;
                                          SIG_S_CURLEN = SIG_S_EncodedSize;
                                          P_STATUS = 1; /*success*/
                                   }
                           }
                     }
             }
      }
       ///*CLEANUP - Release memory structures to prevent memory leaks*/
                                                 /*frees EC KEY memory allocation*/
       EC_KEY_free(eckey);
       EC_GROUP_clear_free(ecgroup);
                                       /*destroys any sensitive data in the group and
then frees the memory*/
      BN clear free(PrivateKeyBN);
                                       /*overwrites the BN before returning memory to
the system*/
       ECDSA_SIG_free(signature);
                                          /*frees ECDSA_SIG memory allocation*/
/*debug logic*/
#if ISaDEBUG
      fp = fopen("debug-siggen.txt", "w");
      fprintf(fp, "hello world\n");
fprintf(fp, " \n");
fprintf(fp, "PrivateKeyBigNumber = ");
      BN_print_fp(fp, PrivateKeyBN);
      fprintf(fp, " \n");
fprintf(fp, "signature->r = ");
      BN_print_fp(fp, signature->r);
      fprintf(fp, " \n");
fprintf(fp, "signature->s = ");
      BN_print_fp(fp, signature->s);
      fprintf(fp, " \n");
      fclose(fp);
#endif
  //END OF CUSTOM CODE
}
```

### A.5 – SigVerify

```
File:
                 sigverify.c
Author:
                 kfischer
Creation date: 12/07/2013 - 18:52
POU name: SigVerify
---autocode---
//START OF CUSTOM CODE HEADERS
#include <debug.h>
#include <stdio.h>
#include <stdlib.h>
#include <openssl/ec.h> // for EC_GROUP_new_by_curve_name, EC_GROUP_free,
EC_KEY_new, EC_KEY_set_group, EC_KEY_generate_key, EC_KEY_free
#include <openssl/ecdsa.h> // for ECDSA_do_sign, ECDSA_do_verify
#include <openssl/obj_mac.h> // for NID_secp384r1
#include <openssl/bn.h>
                              // for BIGNUM conversions
#include <openssl/bio.h>
#include <openssl/evp.h>
#include <openssl/buffer.h>
#include <math.h>
#include <base64.h>
                               // for base64 conversions
//END OF CUSTOM CODE HEADERS
---autocode---
  //START OF CUSTOM CODE
  //--Declarations
  /*debug variables*/
      #if ISaDEBUG
      FILE *fp;
      int i;
      #endif
      /*status variables*/
      #define P_PTR_PUBLIC_KEY_DATA (((char*)VA_ADDRESS(pvBfData,_P_IEC_DATA-
>public key))+2)
      #define P_PTR_HASH_DATA (((char*)VA_ADDRESS(pvBfData,_P_IEC_DATA->hash))+2)
      #define P_PTR_SIG_R_DATA (((char*)VA_ADDRESS(pvBfData,_P_IEC_DATA->sig_r))+2)
      #define P_PTR_SIG_S_DATA (((char*)VA_ADDRESS(pvBfData,_P_IEC_DATA->sig_s))+2)
      const int set_group_success = 1;
      const int gen_success = 1;
      const int set_public_success = 1;
      const int set private success = 1;
      const int verify_success = 1;
      int set_group_status;
      int set_public_status;
      int verify_status;
      unsigned char hash[128];
```

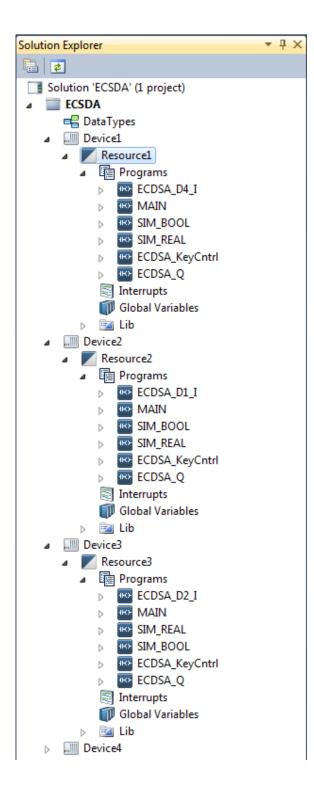
```
/*EC Variables*/
   EC KEY *eckey;
   EC_GROUP *ecgroup;
   EC POINT *PublicKeyPoint;
    ECDSA_SIG *signature;
/*base64 variables for decoding public key*/
   BIGNUM *PublicKeyBN;
   unsigned char *buffer;
 unsigned char *PublicKeyBinaryUnsigned, PublicKeyBinaryUnsignedArray[250];
   const unsigned char *PublicKeyBinaryConst;
    /*base64 variables for decoding SIG R key*/
   BIGNUM *SIG_R_BN;
   unsigned char *SIG_R_buffer;
unsigned char *SIG_R_BinaryUnsigned, SIG_R_BinaryUnsignedArray[250];
   const unsigned char *SIG R BinaryConst;
    /*base64 variables for decoding SIG S key*/
   BIGNUM *SIG S BN;
   unsigned char *SIG_S_buffer;
unsigned char *SIG_S_BinaryUnsigned, SIG_S_BinaryUnsignedArray[250];
   const unsigned char *SIG_S_BinaryConst;
//--Logic
   P_STATUS = 0;
   memcpy(hash,P_PTR_HASH_DATA,128);
   memset(PublicKeyBinaryUnsignedArray,0,250);
   PublicKeyBinaryUnsigned = PublicKeyBinaryUnsignedArray;
   memset(SIG R BinaryUnsignedArray,0,250);
   SIG R BinaryUnsigned = SIG R BinaryUnsignedArray;
   memset(SIG S BinaryUnsignedArray,0,250);
   SIG S BinaryUnsigned = SIG S BinaryUnsignedArray;
   /*STEP 1 - create key object*/
   eckey=EC_KEY_new();
if (NULL == eckey)
 {
     P_STATUS = -1; /*failed to create key object*/
}
else
{
           /*STEP 2 - create EC_GROUP object*/
          ecgroup = EC_GROUP_new_by_curve_name(NID_secp521r1);
     if (NULL == ecgroup)
     {
         P_STATUS = -2; /*failed to create new EC Group*/
    }
    else
     {
                 /*STEP 3 - associate key with group*/
         set group_status = EC_KEY_set_group(eckey,ecgroup);
         if (set group success != set group status)
         {
```

```
P_STATUS = -3; /*failed to associate key with group*/
            }
            else
                     {
                            /*STEP 4 - set public key*/
                            PublicKeyPoint = EC POINT new(ecgroup);
                            PublicKeyBN = BN new();
                            buffer = (unsigned char*) P PTR PUBLIC KEY DATA;
                            EVP DecodeBlock(PublicKeyBinaryUnsigned, buffer, 240);
                            PublicKeyBinaryConst = (const unsigned char*)
PublicKeyBinaryUnsigned;
                            BN_bin2bn(PublicKeyBinaryConst,133,PublicKeyBN);
       EC_POINT_bn2point(ecgroup,PublicKeyBN,PublicKeyPoint,NULL);
                            set_public_status =
EC_KEY_set_public_key(eckey,PublicKeyPoint);
                            if (set_public_success != set_public_status)
                            {
                                   P STATUS = -4; /*failed to set public key*/
                            }
                            else
                            {
                                   /*STEP 5 - convert Base64 SIG_R and SIG_S into BN
and set signature values*/
                                   signature = ECDSA_SIG_new();
                                   SIG_R_BN = signature->r;
                                   SIG S BN = signature->s;
                                   SIG R buffer = (unsigned char*) P PTR SIG R DATA;
       EVP_DecodeBlock(SIG_R_BinaryUnsigned,SIG_R_buffer,120);
                                   SIG_R_BinaryConst = (const unsigned char*)
SIG R BinaryUnsigned;
                                   if (P PTR SIG R DATA[0] == 'A')
                                   {
                                          BN bin2bn(SIG R BinaryConst,66,SIG R BN);
                                   }
                                   else
                                   {
                                          BN_bin2bn(SIG_R_BinaryConst,65,SIG_R_BN);
                                   }
                                   SIG_S_buffer = (unsigned char*) P_PTR_SIG_S_DATA;
       EVP_DecodeBlock(SIG_S_BinaryUnsigned,SIG_S_buffer,120);
                                   SIG_S_BinaryConst = (const unsigned char*)
SIG_S_BinaryUnsigned;
                                   if (P_PTR_SIG_S_DATA[0] == 'A')
                                   {
                                          BN_bin2bn(SIG_S_BinaryConst,66,SIG_S_BN);
                                   }
                                   else
                                   {
                                          BN_bin2bn(SIG_S_BinaryConst,65,SIG_S_BN);
                                   }
                                   if (NULL==signature)
```

```
{
                                           P_STATUS = -5;
                                    }
                                    else
                                    {
                                           /*STEP 6 - Verify Signature*/
                                           verify status =
ECDSA do verify(hash,128,signature,eckey);
                                           if (verify_success != verify_status)
                                           {
                                                   P_STATUS = -6; /*verification failed*/
                                           }
                                           else
                                           {
                                                   P_STATUS = 1; /*signature verified*/
                                           }
                                    }
                            }
                     }
              }
       }
       /*CLEANUP - Release memory structures to prevent memory leaks*/
              EC_KEY_free(eckey);
                                                          /*frees EC_KEY memory
allocation*/
              EC_GROUP_clear_free(ecgroup);
                                              /*destroys any sensitive data in the
group and then frees the memory*/
              EC_POINT_free(PublicKeyPoint); /*frees EC_POINT memory allocation*/
              BN_clear_free(PublicKeyBN); /*overwrites the BN before returning
memory to the system*/
              ECDSA_SIG_free(signature);
                                                  /*frees ECDSA SIG memory allocation*/
/*debug logic*/
#if ISaDEBUG
      fp = fopen("debug-sigverify.txt", "w");
fprintf(fp, "hello world\n");
fprintf(fp, " \n");
fprintf(fp, "PublicKeyBigNumber = ");
       BN_print_fp(fp, PublicKeyBN);
       fprintf(fp, " \n");
fprintf(fp, "signature->r = ");
       BN_print_fp(fp, signature->r);
       fprintf(fp, " \n");
fprintf(fp, "signature->s = ");
       BN_print_fp(fp, signature->s);
       fprintf(fp, " \n");
       fclose(fp);
#endif
}
```

## **APPENDIX B – LADDER LOGIC CODE**

## **B.1 – Solution Explorer**



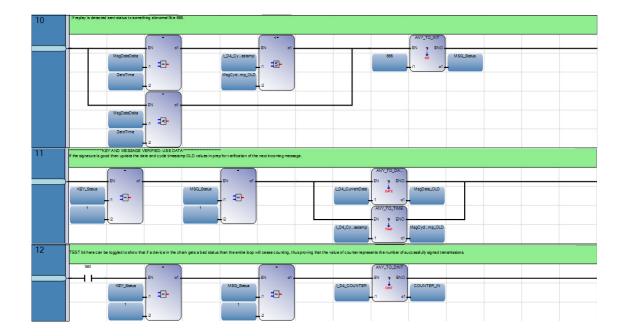
# **B.2 – Bindings Sample**

History of selection			
Producing Groups		1.50 1000 1000	Consuming Groups
Device1	Producing Variables	Consuming Variables	Device1
Resource1	PublicKEY	PublicKEY_D1@ECDSA_D1_I	Resource1
	CB1_CLOSED	I_D1_CB1_CLOSED@ECDSA_D1_I	
Device2     Resource2	CB1_OPENED	I_D1_CB1_OPENED@ECDSA_D1_I	⊡ Device2
	VOLTAGE_A_B	I_D1_VOLTAGE_A_B@ECDSA_D1_I	
Device3	VOLTAGE_A_C	I_D1_VOLTAGE_A_C@ECDSA_D1_I	Device3
ia	VOLTAGE_B_C	I_D1_VOLTAGE_B_C@ECDSA_D1_I	Resource3
Device4	FREQ	I_D1_FREQ@ECDSA_D1_I	Device4
Resource4	COUNTER	I_D1_COUNTER@ECDSA_D1_I	Resource4
	CurrentDate	I_D1_CurrentDate@ECDSA_D1_I	
	SIG_R@ECDSA_Q	I_D1_SIG_R@ECDSA_D1_I	
	SIG_S@ECDSA_Q	I_D1_SIG_S@ECDSA_D1_I	
	KeyGenDate@ECDSA_KeyCntrl	I_D1_KeyGenDate@ECDSA_D1_I	
	KeySIG_R@ECDSA_KeyCntrl	I_D1_KeySIG_R@ECDSA_D1_I	
	KeySIG_S@ECDSA_KeyCntrl	I_D1_KeySIG_S@ECDSA_D1_I	
	CycleTimestamp	I_D1_CycleTimestamp@ECDSA_D1_I	

# B.3 – ECDSA\_D4\_I

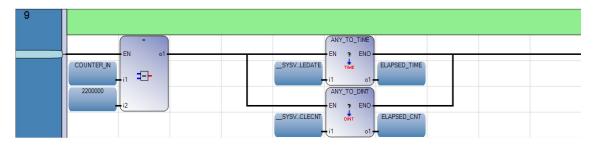
KEY VERIFICATION LOGIC	the key's good.			
	ANY_TO_INT			
EN of PublicKEY_D4	EN PENO	tatus		
, <del>1</del> ⊒-	-11 -1-			
PublickD4_OLD				
If this is a new key, then we need to verify it. This generates the Hash of the new key so we can d	teck to see if it was signed by the old key. Note that TIME is n	t needed here so a constant value of t#0s is used		
		needed here so a consummance of wos a used.		
EN of	KeyMagGen MsgGen EN ENO			
	(LD4_KeyGenDate			
Publick_D4_OLD	HOS			
12	BlankHash			
	HASHJ			
	FALSE			
	FALSE			
	FALSE B2			
	FALSE 83			
	FALSE			
	FALSE B4			
	FALSE BS			
	PALSE			
	FALSE			
	FALSE B7			
	1 88			
	н			
	2			
	3			
	4			
	1.0			
	20			
	R2			
	R3			
	4.0			
	PublicKEY_D4			
Check to see if the new key has been signed by the o	STRIN	hange to prevent replay attacks.		
EN o1	KeySigVenity SigVenity EN ENO	SUB_DATE		
	PublicKD4_OLD	LD4_KeyGenDate KeyG	GenDateDelta	
	KeyMsg.ASH_Q	ADAT1 SUB		
12	HASH	DAT2		
	LD4_KeySIG_S			
If new key was signed by old key, and the timestamps				
		· · · · ·		
EN of	EN of	EN a1		EN 7 ENO
	KeySigV_STATUS	KeyGenDateDelta	KeySigVSTAT	
Publick_D4_OLD		ZeroTime		ANY_TO_ST
-12	-12	- 2	PublicKEY_D4	EN 7 ENO
				ANY_TO_DA
			(_D4_KeyGenDa	
				–↓i1 o1 –

5	If new key was not verified as signe	ed by old key for any reason, then display error co	de and trigger operator alarm (alarm logic n	ot shown, typical PLC to HMI logic use as	umedj.		
	~		( ° )	ANY_TO_INT			
	PublickEY_D4	of KeySigVSTATUS	EN of	KeySigV_STATUS	KEY_Status		
	PublicKD4_OLD		a 12-	( <u> </u>			
	12	attack then set key status to some alarm value (u	i2	in not shown, having DLC to HMI have use			
6	in there appears to be a key replay.	alack henserkey salos o sone alam valoe (c	e, ooo) ana ingger operator a ann (aann iog		assuneo).		
	EN	01	EN of	EN 7 ENO			
	PublickEY_D4	KeyGenDateDelta	a <b>1</b> 3-	666 iii of .	KEY_Status		
	PublicK_D4_OLD	ZeroTime		<u> </u>	-		
7	MSG VERIFICATIO	DN LOGIC	12				
1		ad data so that the signature can be verified.	Merilian 2	Merilian 3			
		ENO	MsgGer_2 MsgGen EN ENO	MsgGer_3 MsgGen EN ENO			
	LD4_CurrentDate	TAT.	DATE STAT	DATE STAT.	-		
	LD4_Cy_estamp	LD4_Cyestamp	TIME HASH	LD4_Cyestamp			
	BlankHash	MsgGenASH_Q		MsgGenASH_Q			
	LD4_C_LOSED	LD4_CLOSED	HASHJ	LD4_CLOSED			
	LD4_CPENED	LD4_CPENED	B1	LD4_CPENED			
	FALSE B2	FALSE	82	FALSE B2			
	- 83		83	= 83			
	TRUE ++ 84	TRUE	84	TRUE			
	FALSE B5	FALSE	95	FALSE			
	TRUE	TRUE		TRUE			
	FALSE B6	FALSE	80	FALSE			
	TRUE B7	TRUE	87	TRUE B7			
		(LD4_COUNTER	88	- 38			
			n	- 1			
	14	14	12	14			
	15	15	13	15			
	16	16		16			
	LD4_V-GE_A_B	LD4_VGE_A_B		LD4_VGE_A.B			
	LD4_V-GE_A_C	LD4_VGE_A_C	R1	LD4_VGE_A_C			
	R2	(L04_VGE_B_C	R2	R2			
	R3	L04_FREQ	R3	R3			
	R4	(D4)FRED	R4	R4			
	'msggen1'	'msggen2'	STRIN_	'msggen3'			
8	Determines status of current messa	ge and signature to determine if the message was	s indeed signed by "PublicKEY_D1".				
	MsgSigVe SigVerit	y nifr	SUB_DATE				
	PublicKEY_04	(_D4_CurrentDate	EN ENO MsgDateDelta				
	PUBLI. S (MsgGenASH_Q	TAT MsgDate_OLD	DAT1 SUB_				
	LD4_SIG_R		DAT2				
	SIG_R						
	LD4_SIG_S						
9	We then check to see that messag	es are timestamped sequentially, or else we might	be really talking to a different device or ad	rotary.			
					ANY_TO_INT		
	MsgDateD	eta			MsgSigVSTATUS	M9G_Status	
	ZeroTim	. 12					
		2					
		EN of	EN of				
	MsgDateD	••••••••••••••••••••••••••••••••••••••	LD4_Cy_estemp				
	ZeroTim		MsgCydmp_OLD				



# **B.4 – MAIN**

	Current Date is obtained and is required for both k	ey generation and transmission, an	nd sending signed data.			
	CURRENT_L		ANY_TO_TIME			
		-D	EN PENO	- 1 -		
		rrentDateSYSVLEDATE	ti1 o1	CycleTimestamp		
2	Due to ISaGRAF limitations on DATE (only goes to updated at intervals greater than 1 second.	seconds) the CycleTimestamp va	riable will also be needed	to protect against message	replay attacks. Key transmission	n will not need this as keeps are
	ANY_TO_TIME					
	EN 2 ENO					
	SYSV_LEDATE TIME Cycle	Timestamp				
3			, <u>,</u>			
						FS
	EN 01					O
	0					
4	- i2					
						ALWAYS_ON
5						Ŭ
	ALWAYS_OFF ALWAYS_OFF					ALWAYS_OFF
						O
6	This rung is just a tool that generates keys which or should only be done during actual commissioning,	an be used to initialize the routine	. The programer will have	to manually copy these initiation of the second sec	al keys into the public and private	e global keys. Note that this
	should only be done during actual commissioning,	and that the FLC will inmediately i	eplace these initialized k	eys with new keys. This run	y can be deleted and replaced v	ith a simple C program if desired.
	FS Ke	vGen_1	eplace these initialized k	<i>KeyVerify_1</i> KeyVerify	g can be deleted and replaced v	ith a simple C program if desired.
	FS		eplace these initialized k		g can be deleted and replaced v	ith a simple C program if desired.
	FS Ke	PUBLL	PublicKEY_Initial	<i>KeyVerify_1</i> KeyVerify	g can be deleted and replaced y	rth a simple C program f desired.
	FS Ke	PUBLI	PublicKEY_Inisal	KeyVerify_7 KeyVerify N ENO PUBLI_ STAT_		ith a simple C program if desired.
	FS Ke	PUBLL	PublicKEY_Inisal	<i>KeyVerify_1</i> KeyVerify N ENO		rth a simple C program if desired.
7	FS Ke	PUBLI	PublicKEY_Inisal	KeyVerify_7 KeyVerify N ENO PUBLI_ STAT_		rth a simple C program if desired.
7		yGen_1 ENO PublicKEY_Inital PUBLI PrivateKEY_Inital PRIVA STAT_	PublicKEY_Initial	KeyVerify_T KeyVerify IN ENO UBLL STAT_ PRIVA_		ith a simple C program if desired.
7	FS ANY	yGen_T ENO PublicKEY_Inital PUBLI PrivateKEY_Inital PRIVA. STAT	eplace these initialized k	KeyVerity_T KeyVerity IN ENO UBLL STAT_ IRIVA_		ith a simple C program if desired.
7	FS ANY	yGen_J ENO PublicKEY_Inital PUBLI PrivateKEY_Inital PRIVA STAT_	eplace these initialized k	KeyVerify NevVerify NUBLL_STAT_ PRIVA_	TER_N	ith a simple C program if desired.
	FS ANY	VGen_T ENO PublicKEY_Initial PUBLL PrivateKEY_Inital PRIVAL STAT	PublicKEY_Initial	KeyVerify KeyVerify N ENO PUBLL STAT- PRIVA- ANY_TO_DINT N 2 ENO DINT COUN		rth a simple C program if desired.
7		VGen_T ENO PublicKEY_Initial PUBLL PrivateKEY_Initial PRIVAL STAT_ TO_DINT ENO COUNTER	PublicKEY_Initial	KeyVerify KeyVerify N ENO PUBLL STAT- PRIVA- ANY_TO_DINT N 2 ENO DINT COUN		ith a simple C program if desired.
		yGer_7 ENO PUBLL PrivateKEY_Initial PRIVA_ STAT_ TO_DINT other other COUNTER	PublicKEY_Initial		TER_N	ith a simple C program if desired.
		VGen_T eNO PublicKEY_Initial PUBLL PrivateKEY_Initial PRIVAL STAT_ TO_DNT ENO COUNTER o1	eplace these initialized k		TER_N ENO.	ith a simple C program if desired.
		yGer_7 ENO PUBLL PrivateKEY_Initial PRIVA_ STAT_ TO_DINT other other COUNTER	eplace these initialized k		TER_N ENO.	ith a simple C program if desired.
		VGen_T eNO PublicKEY_Initial PUBLL PrivateKEY_Initial PRIVAL STAT_ TO_DNT ENO COUNTER o1	eplace these initialized k		TER_N ENO.	ith a simple C program if desired.
		VGen_T eNO PublicKEY_Initial PUBLL PrivateKEY_Initial PRIVAL STAT_ TO_DNT ENO COUNTER o1	PublicKEY_Initial PublicKEY_Initial PrivateKEY_Initial 0		TER_N ENO.	ith a simple C program if desired.



Note: Rungs 8 and 9 are only used for the main routine of Device 1 to initialize the device and to trap the number elapsed time and cycle counts. The other devices have a rung similar to rung 8 but without the FS XIC latch.

# B.5 – SIM\_BOOL

1		
	CB1_CLOSED TOW_1 TON	CB1_CLOSED
		0
2		
	CB1_CLOSED TON 2 TON Q	
3		
	CB1_CLOSED	CB1_OPENED
		<u> </u>

# B.6 – SIM\_REAL

1	
	SIM_TONQ SIM_TON
	TON
	HPT ET
2	
	ANY_TO_RECOS
-	
3	
	NOISE_FACTOR EN ENO EN ENO VOLTAGE_A_B
4	
4	
	NOISE_FACTOR VOLTANOISE 480.0 VOLTAGE_A_C
	10.0 VOLTA_NOISE
	- 12 12
5	VOLTAGE_B_C SIMULATION
	NOISE_FACTOR VOLTANOISE 480.0 VOLTAGE_B_C
	10.0 VOLTA_NOISE
6	FREQ SIMULATION
	NOISE_FACTOR FREQ_NOISE 60.0 FREQ
	1.0 FREQ_NOISE

# B.7 – ECDSA\_KeyCntrl

1		
	START KEY CREATION PROCESS -Controls fraguency of key generation KeyGorOl KayGanTOV KayGanTOV KeyGanTOV	TART
		2
2	Saves OLD keys before generating new keys.	
	Kajŝereŭn         Kajŝereŭn         ANV_TO_ST.         ANV_TO_ST.	eDN
3	Generates new keys and timestamps the new keys.	
	YayGerDN KeyGer_1 KeyGer_1 KeyGer_1 KeyGer_1 KeyGer_1 KeyGer_2 ANY_TO_DA	
	EN ENO PABIOREY PABIOREY Currentizes Dir KeyGenDes	_
	Private/EY Private/EY	
4	STATE	
	KeyderoN KeyderoN Magdan EN ct EN	
		_
	LINE HARL	
	FALSE - B1	
	FALSE	
	FALSE	
	B3	
	FALSE as	
	FALSE	
	FALSE BO	
	e7	
	2	
	u	
	10	
	20	
	R2 82	
	+R3	
	- R4	
	Publicker STRIN.	

	KeyVerif_STATUS		EN ENO	EN ENO	
	_, <del>.</del> ==-	PrivateKE	LOLD Keysig_s		
	1	KeyMsg/		KeyMsg.ASH_Q	
				KeySIG_R	
			STAT.	KeySIQ_S	
Ends KeyGen pro	ooess by confirming new keys were generated a	nd successfully verified.			
	KeyGenDN	•	•		KeyGe
	PublickEY	EN of	PrivateKEY	of EN of KeySigV_STATUS	
		a ==-	, = <del></del>		
	PublicKEY_OLD	12	PrivateKEY_OLD	1	
	KeyGenDN KeyGenSTART				

# B.8 – ECDSA\_Q

START OF OUT	PUT DATA CONTROL						
	MsgGen_1 MsgGen EN ENO	MsgGen_2 MsgGen EN ENO	MsgGen_3 MsgGen EN ENO				
CurrentDate		CurrentDate	CurrentDate				
CycleTimesta	DATE STAT.	DATE STAT.	DATE STAT.				
Cyclemines	TIME HASH	TIME HASH.					
BlankHash		MisgGen., ASH_Q	(M sgGen ASH_Q				
CB1_CLOSE	HASHJ	HASH_J CB1_CLOSED	HASHJ CB1_CLOSED				
	- B1	- B1	- B1				
CB1_OPENE		CB1_OPENED	CB1_OPENED				
FALSE		FALSE	FALSE				
	- B3	- 83	- B3				
TRUE	1.4	TRUE					
FALSE		FALSE	FALSE				
TRUE	<b></b> 85	TRUE B5	TRUE				
TRUE	- 36	- 96					
FALSE		FALSE	FALSE				
TRUE	87	TRUE B7	TRUE B7				
			- 38				
1		COUNTER					
2	-1"	2	2				
	12	12	- 12				
3		3	3				
4	-1"	4					
VOLTAGE_A			VOLTAGE A B				
VOLTAGE_A	→R1						
VOLTAGE_A	c	VOLTAGE_A_C	VOLTAGE_A_C				
VOLTAGE_B	- R2	R2 VOLTAGE_B_C	R2				
		-R3	-R3				
FREQ		FREQ	FREQ				
'msggen1'	-R4	R4 'msggen2'	maggen3				
	STRIN_	STRIN_					
	SigGer_1	SigNenity_1 SigVenity					
}	EN ENO	EN ENO					
PrivateKEY	PRIVA. SIG_R	PublickEY					
MisgGen., ASH							
	HASH SIG_S	HASH SIG_R					
	STAT.	SIG_R					
		SIG_S					