OSTEP Chapter 40

ECE 3600, Fall 2022

Table of Contents

<u>1. File System Implementation</u> <u>2. Inode Contents</u>

2. Indde contents3. Directories4. Open and Read Access Paths5. Create and Write Access Paths6. Exercises

1. File System Implementation

Example with 64 blocks, block size 4 KB:



One block each for superblock, inode bitmap, data bitmap.

5 blocks for inode table (256 bytes per inode, 16 inodes per block), 56 blocks for file data:



Reading an inode: address = 12KB + 256 * inumber; block = address / 4KB; offset = address % 4KB

2. Inode Contents

<u>stat()</u> shows subset of the inode contents; also see <u>inode(7)</u>

mode includes file type

Size	Name	What is this inode field for?
2	mode	can this file be read/written/executed?
2	uid	who owns this file?
4	size	how many bytes are in this file?
4	time	what time was this file last accessed?
4	ctime	what time was this file created?
4	mtime	what time was this file last modified?
4	dtime	what time was this inode deleted?
2	gid	which group does this file belong to?
2	links_count	how many hard links are there to this file?
4	blocks	how many blocks have been allocated to this file?
4	flags	how should ext2 use this inode?
4	osď1	an OS-dependent field
60	block	a set of disk pointers (15 total)
4	generation	file version (used by NFS)
4	file_acl	a new permissions model beyond mode bits
4	dir_acl	called access control lists

Figure 40.1: Simplified Ext2 Inode

Multi-level indexing for larger files: indirect block pointers, double indirect, triple indirect

3. Directories

Directories are just files with a special structure.

inum	reclen	strlen	name
5	12	2	
2	12	3	
12	12	4	foo
13	12	4	bar
24	36	28	foobar_is_a_pretty_longname

4. Open and Read Access Paths

open("/foo/bar", 0_RDONLY)

	data	inode								
	bitmap	bitmap	inode	inode	inode	data	data	data [0]	data [1]	data [2]
			read							
						read				
open(bar)				read						
							read			
					read					
					read					
read()								read		
					write					
read()					read					
									read	
					write					
read()					read					
										read
					write					

 Figure 40.3: File Read Timeline (Time Increasing Downward)

5. Create and Write Access Paths

	data bitmap	inode bitmap		foo inode			foo data	bar data [0]	bar data [1]	bar data [2]
			read	read		read				
create (/foo/bar)		read write					read write			
				write	read write					
write()	read write				read					
					write			write		
write()	read write				read				write	
					write				witte	
write()	read write				read					write
					write					

Figure 40.4: File Creation Timeline (Time Increasing Downward)

6. Exercises

Exercises from the book using <u>vsfs.py</u>:

\$ python ./vsfs.py -n 4

Initial state

inode bitmap inodes [d a:0 r:2] [] [] [] [] [] [] [] data bitmap data [(.,0) (..,0)] [] [] [] [] [] [] [] []

Which operation took place?

inode bitmap 11000000 inodes [d a:0 r:3] [d a:1 r:2] [] [] [] [] [] [] data bitmap 11000000 data [(.,0) (..,0) (g,1)] [(.,1) (..,0)] [] [] [] [] [] []

Which operation took place?

inode bitmap	11100000
inodes	[d a:0 r:3] [d a:1 r:2] [f a:-1 r:1] [] [] [] [] []
data bitmap	11000000
data	[(.,0) (,0) (g,1) (q,2)] [(.,1) (,0)] [] [] [] [] [] []

Which operation took place?

inode bitmap	11110000
inodes	[d a:0 r:3] [d a:1 r:2] [f a:-1 r:1] [f a:-1 r:1] [] [] [] []
data bitmap	11000000
data	[(.,0) (,0) (g,1) (q,2) (u,3)] [(.,1) (,0)] [] [] [] [] [] []

Which operation took place?

inode bitmap inodes [d a:0 r:3] [d a:1 r:2] [f a:-1 r:1] [f a:-1 r:2] [] [] [] data bitmap data [(.,0) (..,0) (g,1) (q,2) (u,3) (x,3)] [(.,1) (..,0)] [] [] [] [] [] []

\$ python ./vsfs.py -n 4 -r

Initial state

inode bitmap inodes [d a:0 r:2] [] [] [] [] [] [] [] data bitmap data [(.,0) (..,0)] [] [] [] [] [] [] []

mkdir("/g"); State of file system (inode bitmap, inodes, data bitmap, data)?

creat("/q");

creat("/u");

link("/u", "/x");